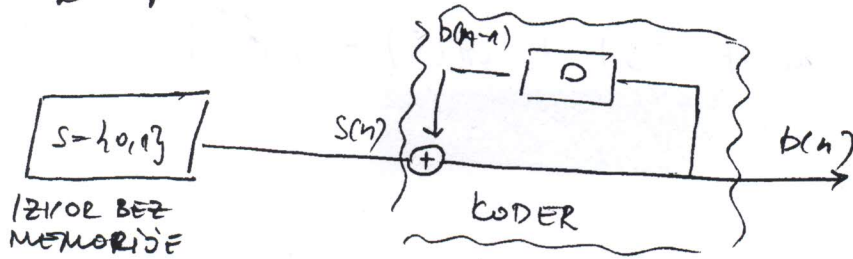


7. Binarni izvor bez memorije sa verovatnoćama simbole $P(1^n) = p$ i $P(0^n) = q = 1-p$ dovodi se na ulaz kodera kao na slici:



Odrediti entropiju binarnog niza na izlazu kodera sa slike.

Rešenje:

$$b(n) = s(n) \oplus b(n-1)$$

$$\left. \begin{array}{l} s(n) = \{00100110000100 \\ b(n) = \{0\}0011101111000 \end{array} \right\} \Rightarrow \text{diferencijalni kod}$$

$b(-1) = \text{početno stanje}$

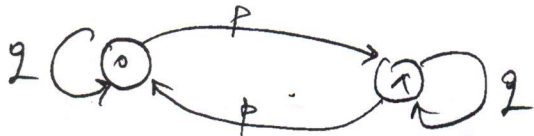
odgovarajućo izvor je sa memorijom 1 reda:

$$P(b(n)=0 | b(n-1)=0) = P(0/0) = P(s(n)=0) = q$$

$$P(b(n)=1 | b(n-1)=0) = P(1/0) = P(s(n)=1) = p$$

$$P(b(n)=0 | b(n-1)=1) = P(0/1) = P(s(n)=1) = p$$

$$P(b(n)=1 | b(n-1)=1) = P(1/1) = P(s(n)=0) = q$$



$$P(0) = P(b(n)=0) = P(b(n)=0) \cdot q + P(b(n)=1) \cdot p = qP(0) + pP(1)$$

$$\left. \begin{array}{l} (1) \quad p \cdot P(1) = (1-q) \cdot P(0) \\ (2) \quad P(0) + P(1) = 1 \end{array} \right\} \Rightarrow P(1) = P(0) = 0.5 \Rightarrow H(B) = \sum_{i=1}^2 P(b_i) \log \frac{1}{P(b_i)} = 1$$

$$H^{(n)}(B) = \sum_{B^{n+1}} P(b_1, b_2, \dots, b_n, b_{n+1}) \log \frac{1}{P(b_j | b_1, b_2, \dots, b_n)}$$

$$H^{(1)}(B) = \sum_{i=1}^2 \sum_{j=1}^2 P(b_j | b_i) P(b_i) \log \frac{1}{P(b_j | b_i)}$$

$$H^{(1)}(B) = P(0/0) P(0) \log \frac{1}{P(0/0)} + P(1/0) P(0) \log \frac{1}{P(1/0)} + P(0/1) P(1) \log \frac{1}{P(0/1)} + P(1/1) P(1) \log \frac{1}{P(1/1)}$$

$$H^{(1)}(B) = 2/2 \log \frac{1}{q} + 1/2 \log \frac{1}{p} + 1/2 \log \frac{1}{p} + 2/2 \log \frac{1}{q}$$

$$H^{(1)}(B) = 1 \log \frac{1}{p} + 2 \log \frac{1}{q} = H(S)$$

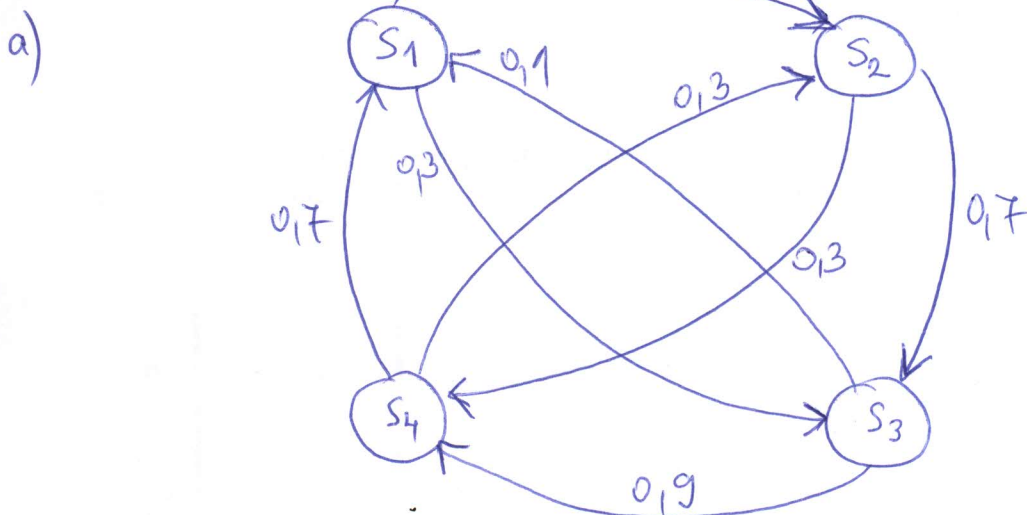
IZVOR SA MEMORIOM PRVOG REDA KOJI EMITUJE SIMBOLE s_1, s_2, s_3 I s_4
 OPISAN JE TRANZICIONOM MATRICOM

$$\Pi = \left[\begin{array}{cccc} 0 & 0,7 & 0,3 & 0 \\ 0 & 0 & 0,7 & 0,3 \\ 0,1 & 0 & 0 & 0,9 \\ 0,7 & 0,3 & 0 & 0 \end{array} \right] \left. \vphantom{\Pi} \right\} s_i \quad \Pi = P(s_j | s_i)$$

s_j

- a) NAERTATI DIAGRAM STANJA KOJI ODGOVARA OVOM IZVORU.
 ODREDITI VEROVATNOĆE PRIDRUŽENOG IZVORA (VEROVATNOĆE STANJA)
 I ENTROPIJU PRIDRUŽENOG IZVORA.
- b) ODREDITI ENTROPIJU IZVORA SA MEMORIOM.

Rešenje:



$$\left. \begin{array}{l} P(s_1) = 0,1 \cdot P(s_3) + 0,7 P(s_4) \\ P(s_2) = 0,7 P(s_1) + 0,3 P(s_4) \\ P(s_3) = 0,3 P(s_1) + 0,7 P(s_2) \\ P(s_4) = 0,3 P(s_2) + 0,9 P(s_3) \end{array} \right\}$$

JEDNU OD JEDNAČINA IZBRACITI
 JER JE LINEARNO ZAVISNA, A
 UMESTO NJE KORISTITI:

$$\sum_{i=1}^4 P(s_i) = 1$$

Rešenje ovog sistema je:

$$\begin{aligned} P(s_1) &= 0,23 \\ P(s_2) &= 0,25 \\ P(s_3) &= 0,24 \\ P(s_4) &= 0,29 \end{aligned}$$

$$H(\bar{S}) = \sum_{i=1}^4 P(s_i) \log_2 \frac{1}{P(s_i)} = 1,994 \text{ sh/simbol}$$

$$b) H(s) = \sum_{i=1}^4 \sum_{j=1}^4 P(s_i) P(s_j | s_i) \cdot \text{ld} \frac{1}{P(s_j | s_i)}$$

$$\begin{aligned} H(s) &= P(s_1) \left(0,7 \text{ld} \frac{1}{0,7} + 0,3 \text{ld} \frac{1}{0,3} \right) \\ &+ P(s_2) \left(0,7 \text{ld} \frac{1}{0,7} + 0,3 \text{ld} \frac{1}{0,3} \right) \\ &+ P(s_3) \left(0,1 \text{ld} \frac{1}{0,1} + 0,9 \text{ld} \frac{1}{0,9} \right) \\ &+ P(s_4) \left(0,7 \text{ld} \frac{1}{0,7} + 0,3 \text{ld} \frac{1}{0,3} \right) \end{aligned}$$

$$H(s) = 0,7825 \text{ sh/symbol}$$

$$\Delta H(s) = H(\bar{s}) - H(s) = 1,2112 \text{ sh/symbol}$$

$$\boxed{H(\bar{s}) > H(s)} !$$

5) Nacetati dijagram stanja Markovljevskog izvora II reda sa binarnom listom simbola $s = \{0, 1\}$ i sa uslovnim verovatnoćama pojavljivanja simbola:

$$P(0|00) = 0,8$$

$$P(1|11) = 0,7$$

$$P(0|00) = 0,2$$

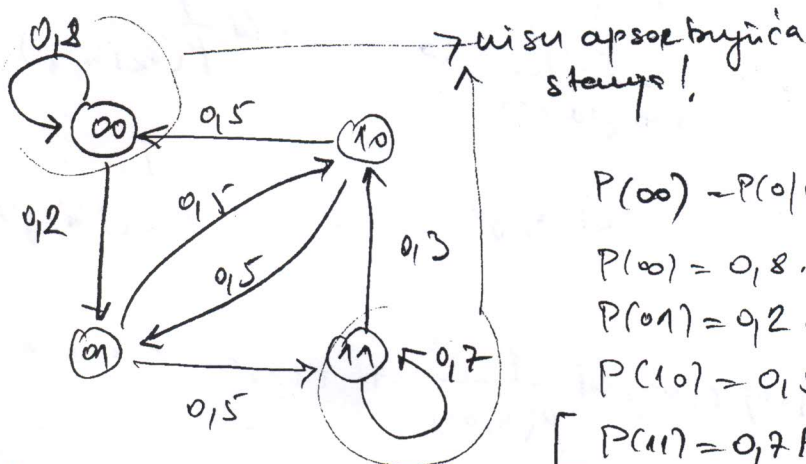
$$P(0|11) = 0,3$$

$$P(0|01) = P(1|01) = P(0|10) = P(1|10) = 0,5$$

ENTROPIJA
II REDA
PROBAJ!

Odeediti verovatnoće stanja pridenenog izvora i verovatnoće izlaznih simbola.

Rešenje:



$$P(00) = P(0|00)P(00) + P(0|10)P(10)$$

$$P(00) = 0,8 \cdot P(00) + 0,5 \cdot P(10) \quad (1)$$

$$P(01) = 0,2 \cdot P(00) + 0,5 \cdot P(10) \quad (2)$$

$$P(10) = 0,5 \cdot P(01) + 0,3 \cdot P(11) \quad (3)$$

$$[P(11) = 0,7 \cdot P(11) + 0,5 \cdot P(01)] \quad X$$

$$P(00) + P(01) + P(10) + P(11) = 1 \quad (4)$$

(suma prethodne četiri jednačine)

$$(1) \Rightarrow 0,2 P(00) = 0,5 P(10) \Rightarrow P(00) = 2,5 P(10)$$

$$(2) \Rightarrow P(01) = 0,5 P(10) + 0,5 P(10) = P(10) \Rightarrow P(01) = P(10)$$

$$(3) \Rightarrow 0,5 P(10) = 0,3 P(11) \Rightarrow P(11) = \frac{5}{3} P(10)$$

$$(4) \Rightarrow 2,5 P(10) + P(10) + P(10) + \frac{5}{3} P(10) = 1 \Rightarrow \boxed{P(10) = \frac{6}{37}}$$

$$P(00) = \frac{5}{2} \cdot P(10) = \frac{15}{37}$$

$$P(01) = P(10) = \frac{6}{37}$$

$$P(11) = \frac{5}{3} P(10) = \frac{5}{3} \cdot \frac{6}{37} = \frac{10}{37}$$

$$P(0) = P(0|00)P(00) + P(0|01)P(01) + P(0|10)P(10) + P(0|11)P(11)$$

$$P(0) = 0,8 \cdot \frac{15}{37} + 0,5 \cdot \frac{6}{37} + 0,5 \cdot \frac{6}{37} + 0,3 \cdot \frac{10}{37} = \frac{21}{37}$$

$$P(1) = 1 - P(0) = \frac{16}{37} \quad \left[\text{tj. } P(1) = P(00) + \frac{1}{2} [P(01) + P(10)] \right]$$

Pridruženi izvor (\bar{S})

000 SU STANJA, A NE SIMBOLI IZVORA!

| Stavec | P_i |
|--------|-------|
| 00 | 15/37 |
| 01 | 6/37 |
| 10 | 6/37 |
| 11 | 10/37 |

PAZI!

$$H(\bar{S}) = \sum_{i=1}^n P_i \log \frac{1}{P_i} = 1,8894 \frac{\text{sh}}{\text{simb}}$$

$s_1=0, s_2=1 \Rightarrow \bar{S} = \begin{cases} P(0) = \frac{21}{37} \\ P(1) = \frac{16}{37} \end{cases} \Rightarrow H(\bar{S}) = 0,9751 \frac{\text{sh}}{\text{simb}}$

$$H(S) = \sum_{s_1, \dots, s_n} P(s_1, s_2, \dots, s_n, s_j) \log \frac{1}{P(s_j | s_1, s_2, \dots, s_n)}$$

$$H(mS) = H(2S) = \sum_{i=0}^1 \sum_{j=0}^1 \sum_{k=0}^1 P(s_i, s_j, s_k) \log \frac{1}{P(s_k | s_i, s_j)}$$

$$H(2S) = \sum_{i=0}^1 \sum_{j=0}^1 \sum_{k=0}^1 P(s_k | s_i, s_j) P(s_i, s_j) \log \frac{1}{P(s_k | s_i, s_j)}$$

$$H(2S) = P(0|00)P(00) \log \frac{1}{P(0|00)} + P(1|00)P(00) \log \frac{1}{P(1|00)} +$$

$$+ P(0|01)P(01) \log \frac{1}{P(0|01)} + P(1|01)P(01) \log \frac{1}{P(1|01)}$$

$$+ P(0|10)P(10) \log \frac{1}{P(0|10)} + P(1|10)P(10) \log \frac{1}{P(1|10)}$$

$$+ P(0|11)P(11) \log \frac{1}{P(0|11)} + P(1|11)P(11) \log \frac{1}{P(1|11)}$$

$$H(2S) = 0,8552 \frac{\text{sh}}{\text{simb}}$$

$$\Delta H = H(\bar{S}) - H(S) \approx 0,12 \frac{\text{sh}}{\text{simb}}$$

Ali bi se kodovale Huffmanovo (n-to prostave) moglo bi da se pride entropiji

$$H(S) \Rightarrow \lim_{n \rightarrow \infty} \frac{L_n}{n} \approx H(\bar{S}) \text{ ali } \underline{\text{ni}} \underline{\text{ne}} \underline{\text{d}}$$

tj. ostala bi ΔH neiskoriscena !!!